# kiltera: a language for concurrent, interacting, mobile, timed systems

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### Outline

Introduction

Overview of kiltera processes



• kiltera: a formalism for describing systems which are:





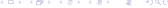
- kiltera: a formalism for describing systems which are:
  - dynamic





- kiltera: a formalism for describing systems which are:
  - dynamic
  - concurrent





- kiltera: a formalism for describing systems which are:
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- kiltera: a formalism for describing systems which are:
  - dynamic
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  - discrete-event



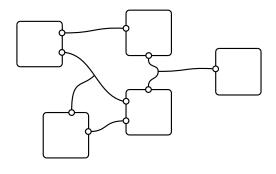


- kiltera: a formalism for describing systems which are:
  - dynamic
  - concurrent
  - interactive
  - timed
  - discrete-event
  - mobile



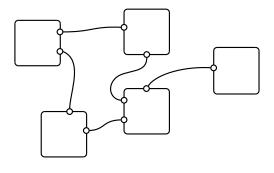


#### A network of processes:



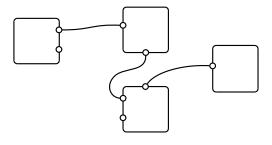


### Mobility: channels can move



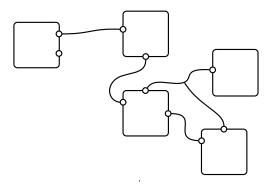


Mobility: components can disapear



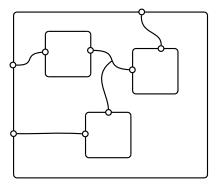


Mobility: components can be created





Nesting: a network of processes is a process





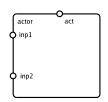
- Time
- Global clock
- Time-base: real numbers
- Time consistency
- Processes can observe and measure the passage of time between events
- Process behaviour can be determined by the timing of events
- Processes can schedule events in the future





Process definitions

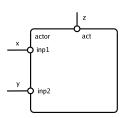
```
process actor[inp1, inp2, act]:
     <body>
```



This defines a class of processes with a given interface.

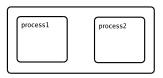


Process instantiation





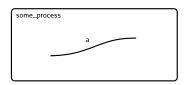
Parallel processes





Channel declarations

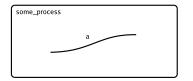
channel a in
 <some process>





#### Channel declarations

async channel a in
 <some process>

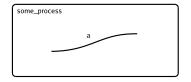


Creates a channel for asynchronous communication (unbounded buffer.) This is the default.



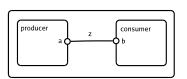
Channel declarations

sync channel a in
 <some process>

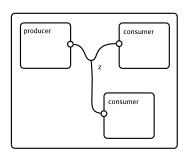


Creates a channel for synchronous communication (rendez-vous.)



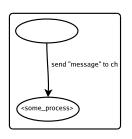








Sending messages

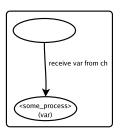


Sending is blocking if the channel is synchronous.



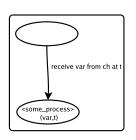
Receiving messages

receive var from ch ->
 <some\_process>





receive var from ch at t ->
 <some process>



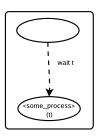
The variable t becomes bound to the time elapsed since the process arrived at the receiving state, until it synchronizes.





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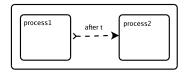
# Overview of kiltera processes Waiting





#### **Timeout**

```
timeout
  css1>
after t ->
  cess2>
```

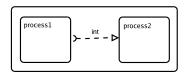


The timeout will take effect if process1> has not finished or has not engaged in any external interaction before time units.



#### Interrupt

```
do
     process1>
interrupt
     process2>
```



coress1> will be interrupted if coress2> engages in
an external interaction.



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#### Other constructs:

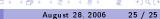
- Data structures
- Conditionals
- Local variable declarations
- Assignment
- Pattern-matching
- Loops
- Recursion
- Channel arrays
- Process arrays
- Auxiliary function definitions





- Language features:
  - Full mobility: moving "live" processes; ambients
  - Multi-way synchronization
  - Deterministic parallel composition





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- Code generation:
  - kiltera2CSP
  - kiltera2C
  - kiltera2DEVS



