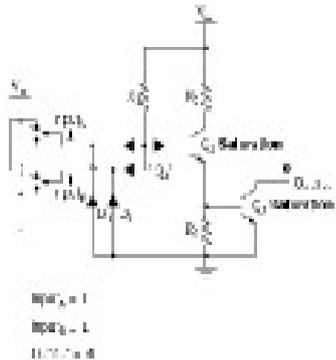


Computer Systemen en Computer Architectuur



Dinsdag 31 oktober 2023 – getalvoorstellingen ((mee)looples)

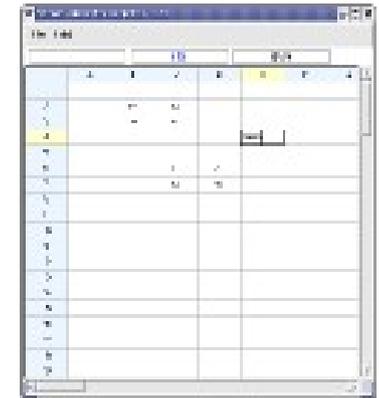
Hans Vangheluwe



```

strcpy:
    addi $sp, $sp, -4
    sw $s0, 0($sp)
    add $s0, $zero, $zero
L1: add $t1, $s0, $a1
    lbu $t2, 0($t1)
    add $t3, $s0, $a0
    sb $t2, 0($t3)
    beq $t2, $zero, L2
    addi $s0, $s0, 1
    j L1
L2: lw $s0, 0($sp)
    addi $sp, $sp, 4
    jr $ra
    
```

$k = 1$
 $xPowerK = x$
 $Sign = 1; s = 0$
 while $k \leq N$:
 $term = sign * xPowerK / factorial(k)$
 $s = s + term$
 $k = k + 2$
 $xPowerK = xPowerK * xSquare$
 $sign = -sign$



physics digital

electronics

computer

architecture / systems

HL progr.
languages
operating
systems

complex SW
applications

Systemen/Engineering

Software Ontwikkeling

Theoretische Informatica

Wiskunde

Jaar 1

Computer Systemen & Architectuur

Project Software Engineering

Computer Graphics

Inleiding Programmeren

Talen & Automaten

Gegevens-abstractie & Data Structuren

Discrete Wiskunde

Jaar 2

Netwerken

Operating Systems

Programming Project Databases

Advanced Programming

Compilers

Inleiding Databases

Algoritmes en Complexiteit

Machines & Berekenbaarheid

Elementaire Statistiek

Lineaire Algebra

Numerieke Analyse

Jaar 3

Distributed Systems

Datastructuren en Graafalgoritmes

Numerieke Lineaire Algebra

Oefeningen ... al doende ...

The image shows a screenshot of the MARS MIPS assembler interface and a Logisim circuit diagram. The MARS window displays assembly code for a Fibonacci program, with the following code visible:

```
6: la $s0, fibs # load address ...
ori $t6,$s0,0 # load address ...
7: la $s5, size # load address ...
ori $t6,$s5,0($s5) # load address ...
8: lw $s5, 0($s5) # load array size
9: li $s2, 1 # 1 is the know...
10: sw $s2, 0($s0) # F[0] = 1
11: sw $s2, 4($s0) # F[1] = F[0] = 1
12: addi $s1, $s5, -2 # Counter for ...
15: loop: lw $s3, 0($s0) # Get value fr...
16: lw $s4, 4($s0) # Get value fr...
17: add $s2, $s3, $s4 # F[n] = F[n-1...
18: sw $s2, 8($s0) # Store newly ...
19: addi $s0, $s0, 4 # increment ad...
20: addi $s1, $s1, -1 # decrement lo...
21: bgtz $s1, loop # repeat while...
```

The Logisim circuit diagram shows a 5-bit shift register labeled 'f5' with 'D', 'Q', and 'clr' inputs. The 'D' input is connected to a box containing the binary value '1.1.1.1.0.1.0.1'. The 'Q' output is connected to a box containing the binary value '1.1.1.1.0.1.0.1'. The 'clr' input is connected to a green clock signal source labeled 'clock'.

MARS

<http://courses.missouristate.edu/KenVollmar/MARS/>

Logisim
<https://sourceforge.net/projects/circuit/>

Data Representation

(binary) Data Representation

- Numbers

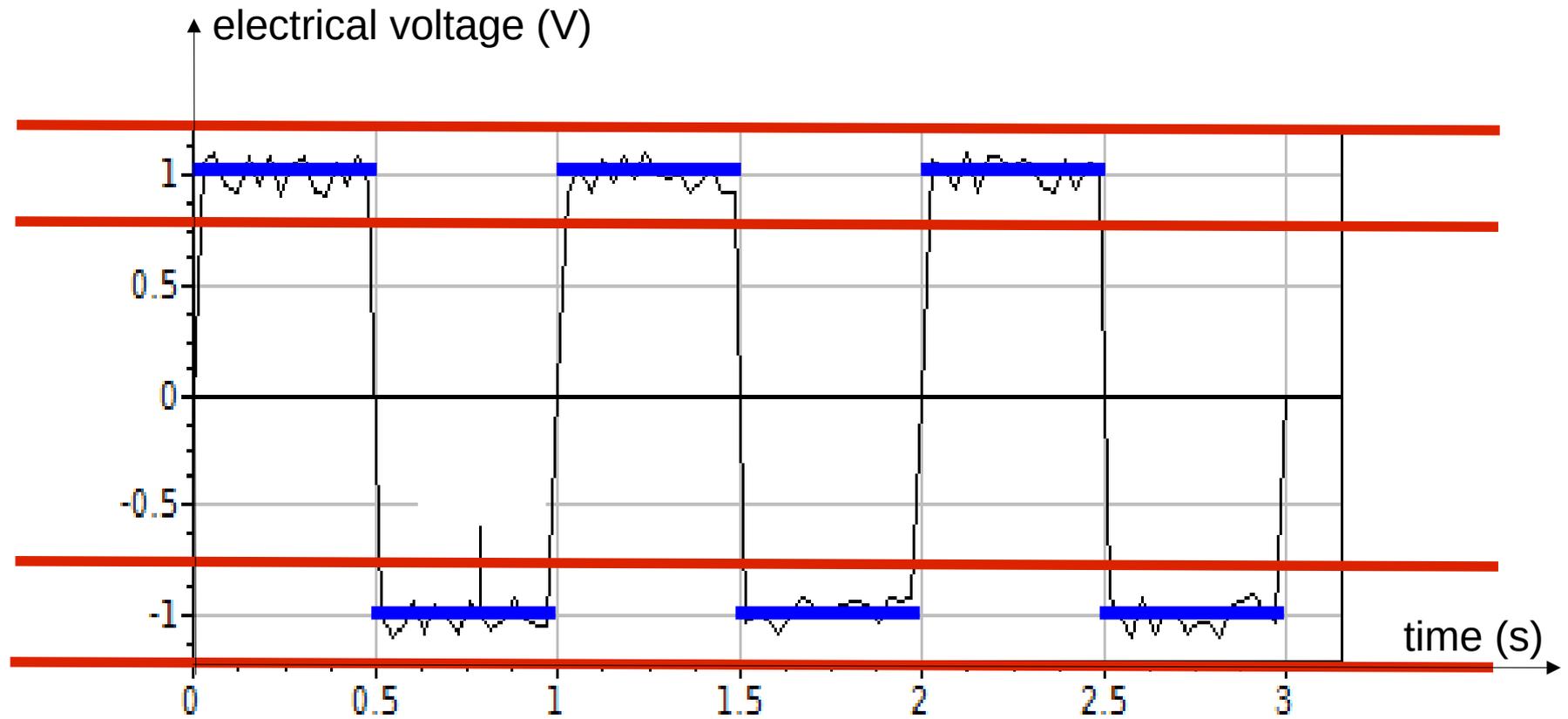
$\mathbb{N}, \mathbb{Z}, \mathbb{Q}, \mathbb{R}$

- Letters (and “strings”)

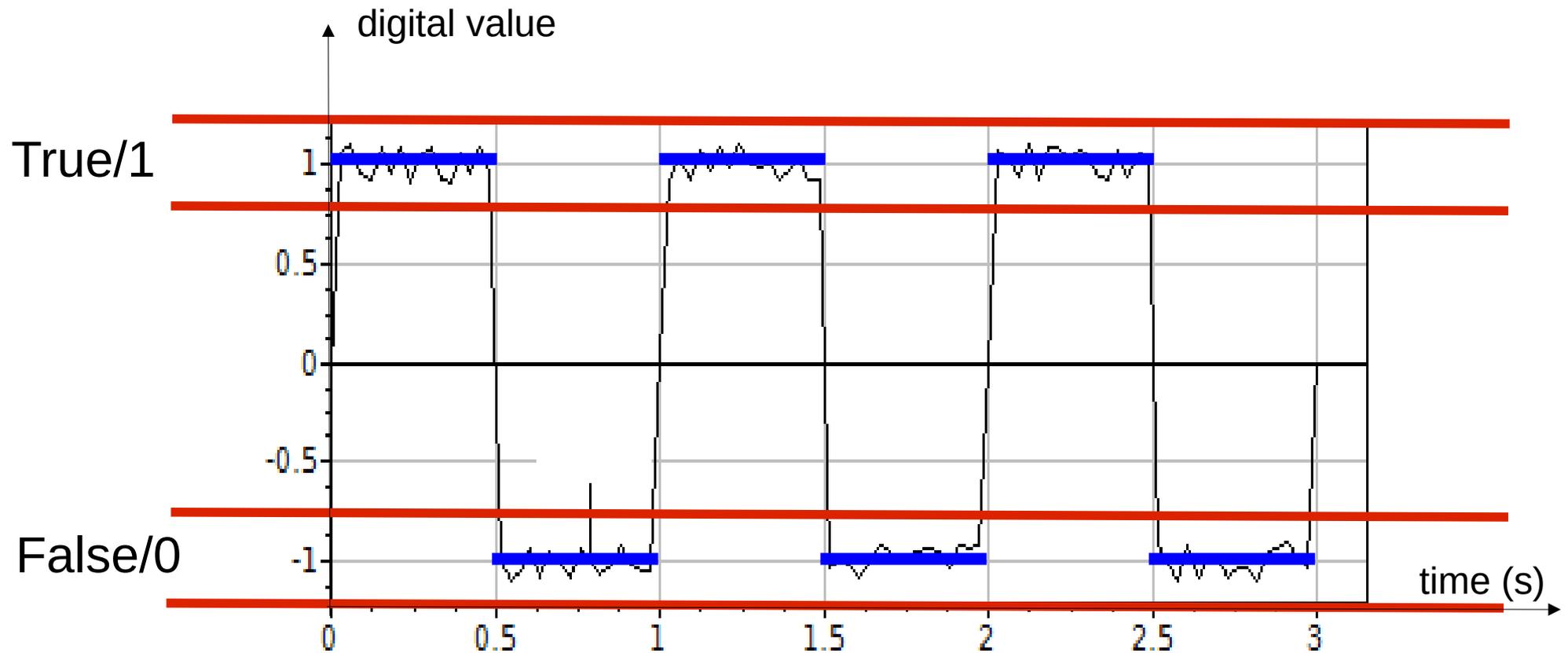
- Instructions

-

from Analog ...



... to Digital



logical 0/1 – Boolean True/False – asserted/de-asserted – high/low

Binary representation

- in computing and telecommunications
- a **bit** is a basic *unit of information storage*
- “**binary digit**”
- the maximum amount of **information** that can be stored in only two distinct states

0 or 1

Bit sequences

(bit) “string” (vs. char string)

0	1 bit	bit	
0111	4 bits	nibble	
01100001	8 bits	byte	French: “octet”
0101010110110111	16 bits	half-word	
...	32 bits	word	
...	64 bits	word	

A **word** is a natural unit of data used by a particular processor design
(historically: 8, 16, 24, 32, or 64 bits)

Bit is abbreviated as “b”. e.g., **kbps** (kilo bits per second)

Byte is abbreviated as “B”. e.g., **kB** (kilo Bytes)

Binary representation ++

With 1 bit, can **represent 2 distinct entities**

With 2 bits, can represent 4 distinct entities

...

With **N** bits, can represent **2^N** distinct entities

Example: {Red, Green, Blue}
encoded as {00, 01, 10}

1 Bit	2 Bits	3 Bits	4 Bits	5 Bits
0	00	000	0000	00000
1	01	001	0001	00001
	10	010	0010	00010
	11	011	0011	00011
		100	0100	00100
		101	0101	00101
		110	0110	00110
		111	0111	00111
			1000	01000
			1001	01001
			1010	01010
			1011	01011
			1100	01100
			1101	01101
			1110	01110
			1111	01111
				10000
				10001
				10010
				10011
				10100
				10101
				10110
				10111
				11000
				11001
				11010
				11011
				11100
				11101
				11110
				11111

Binary representation/approximation of numbers

Binary representation/encoding of **Unsigned Integers** (\mathbb{N})

Binary Coded Decimal (BCD) representation/encoding of **Unsigned Integer** (\mathbb{N}) / **Real** numbers (\mathbb{R})

Binary representation/encoding of **Signed Integer Numbers** (\mathbb{Z})

- Signed magnitude
- One's complement
- Two's complement
- Biased (Excess B)

Fixed-Point binary **approximation**/representation (\mathbb{Q})
of **Real** numbers (\mathbb{R})

Floating Point binary **approximation**/representation
of **Real** numbers (\mathbb{R})

Binary representation/encoding of Unsigned Integers (\mathbb{N})

" $x_{n-1}x_{n-2} \dots x_1x_0$ " is an n-bit, **base 2** (binary) encoding of value x

$$x = x_{n-1}2^{n-1} + x_{n-2}2^{n-2} + \dots + x_12^1 + x_02^0$$

x_0 Least Significant Bit (**LSB**)

x_{n-1} Most Significant Bit (**MSB**)

■ Range: $[0, +2^n - 1] \subseteq \mathbb{N}_0$

■ Example

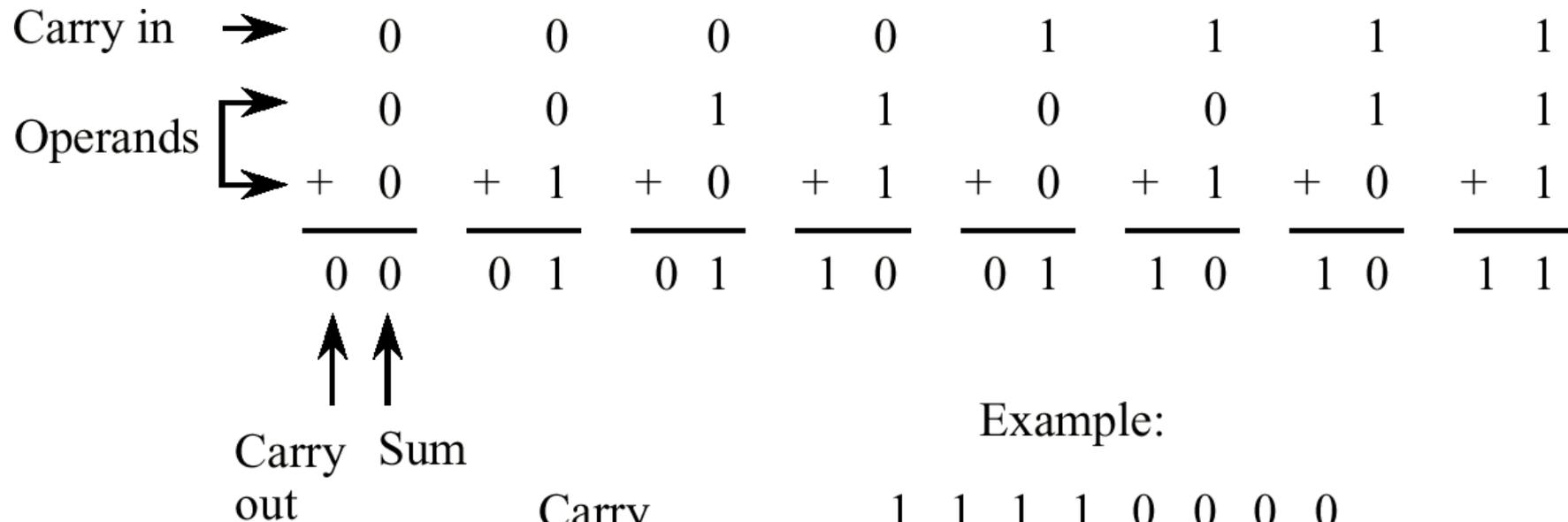
$$0000 \ 0000 \ 0000 \ 0000 \ 0000 \ 0000 \ 0000 \ 1011_2$$

$$= 0 + \dots + 1 \times 2^3 + 0 \times 2^2 + 1 \times 2^1 + 1 \times 2^0$$

$$= 0 + \dots + 8 + 0 + 2 + 1 = 11_{10}$$

1010001001010001010
111011001011001010101
0101 THERE ARE 101
001011 ONLY 10 10010
10010 TYPES OF 0010
101110 PEOPLE: 01011
0010 THOSE WHO 000
100 UNDERSTAND 011
110101 BINARY 10101
0001 AND THOSE 000
1010 WHO DON'T 000
101011101100110101010
11001010101010101000
1010010010100101000
101011101100110101010

(intermezzo) Binary addition of Unsigned Integers (\mathbb{N})



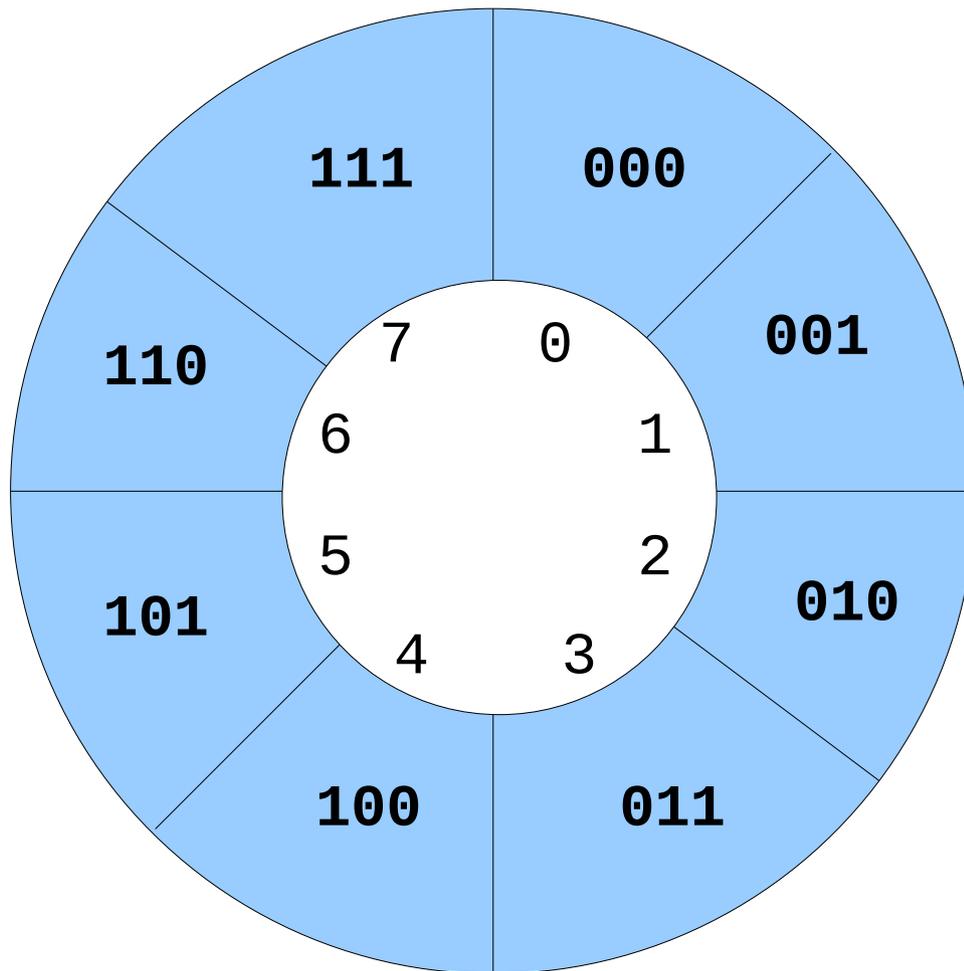
Example:

Carry	1	1	1	1	0	0	0	0		
Addend: A	0	1	1	1	1	1	0	0	$(124)_{10}$	
Augend: B	+	0	1	0	1	1	0	1	0	$(90)_{10}$
Sum									$(214)_{10}$	
	1	1	0	1	0	1	1	0		

unsigned "modulo" arithmetic

$$X = x_{n-1}2^{n-1} + x_{n-2}2^{n-2} + \dots + x_12^1 + x_02^0$$

+1 (unsigned): carry gets discarded, no "overflow"
→



000	0
001	1
010	2
011	3
100	4
101	5
110	6
111	7

$$\begin{aligned} N \text{ bits: } [[2^N]] &= (2^N) \text{ MOD } 2^N = 0; \\ & (2^N) \text{ DIV } 2^N = 1 \\ [[2^{N+k}]] &= (2^{N+k}) \text{ MOD } 2^N = k \end{aligned}$$

used in for example address calculation (PC)

Binary Coded Decimal (BCD)

representation of unsigned Int (\mathbb{N}) / Real (\mathbb{R})

Decimal:	0	1	2	3	4	5	6	7	8	9
BCD:	0000	0001	0010	0011	0100	0101	0110	0111	1000	1001

Every Decimal digit gets represented by 4 bits (Binary digits)

Decimal:	127	:	1	2	7
BCD:	000100100111	:	0001	0010	0111

Used mostly in

- Mainframes/servers (financial applications)
- Embedded microcontrollers/small processors (computation <<<)
- Only digital logic (no processor), display (7-segment)

Word size often 10 or 12 decimal digits (e.g., in calculators)

Binary Coded Decimal (BCD)

Advantages:

- Easy to print, display (e.g., 7-segment), ... thanks to **per-digit** conversion
- Numbers such as decimal 0.2 have
an **infinite** representation in binary
(0.001100110011...)
a **finite** representation in binary-coded decimal
(0000.0010)
- Scaling by factors of 10 by **shifting** (clever compiler ...)
- 10-based **rounding** is easy

Disadvantages:

- More **complex** to implement +, -, *, / (+: 15-20% more circuitry)
- Slower
- More storage space required
e.g., 15_{10} : 1111_{binary} vs. 00010101_{BCD}