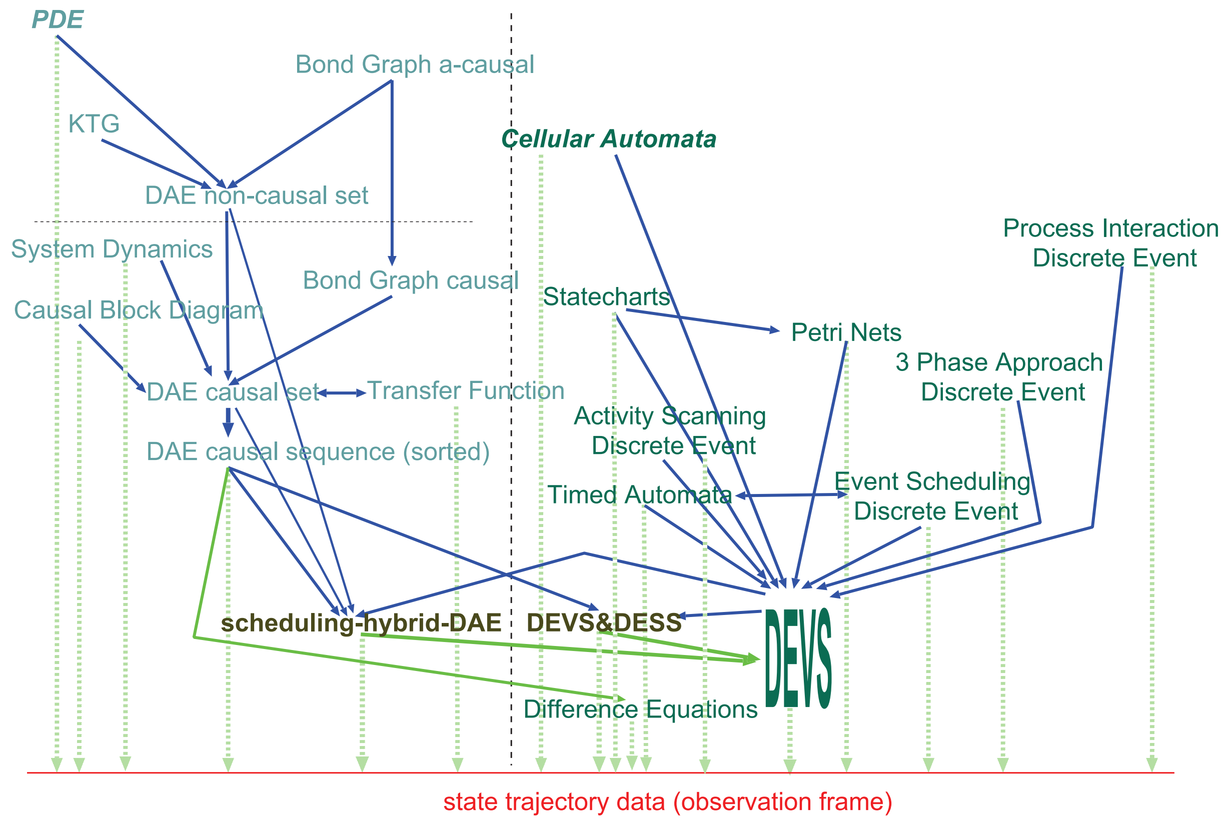
Discrete EVent System specification (DEVS)

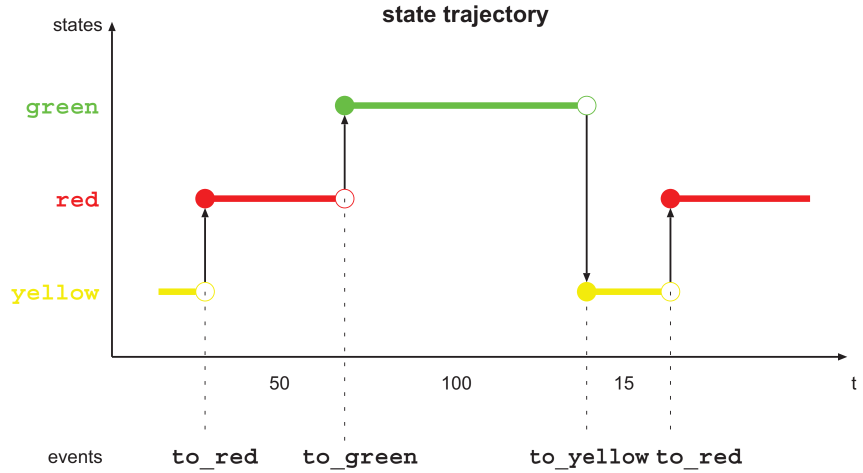
Bernard Zeigler (1976 "Theory of Modelling and Simulation")

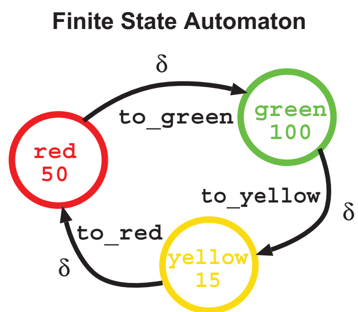
- A formal basis
- for (low-level) representation
- of *all* discrete event modelling formalisms
  (and even others, after approximation)
- and simulator implementations

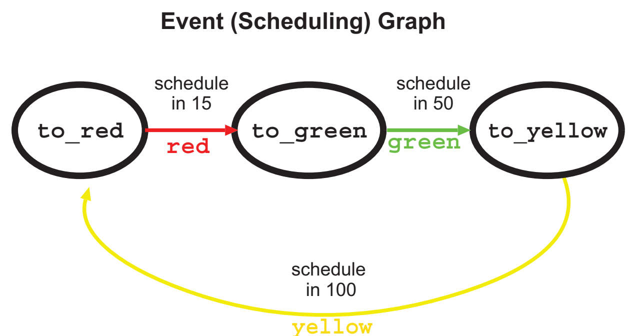
DEVS's central place in the Formalism Transformation Graph



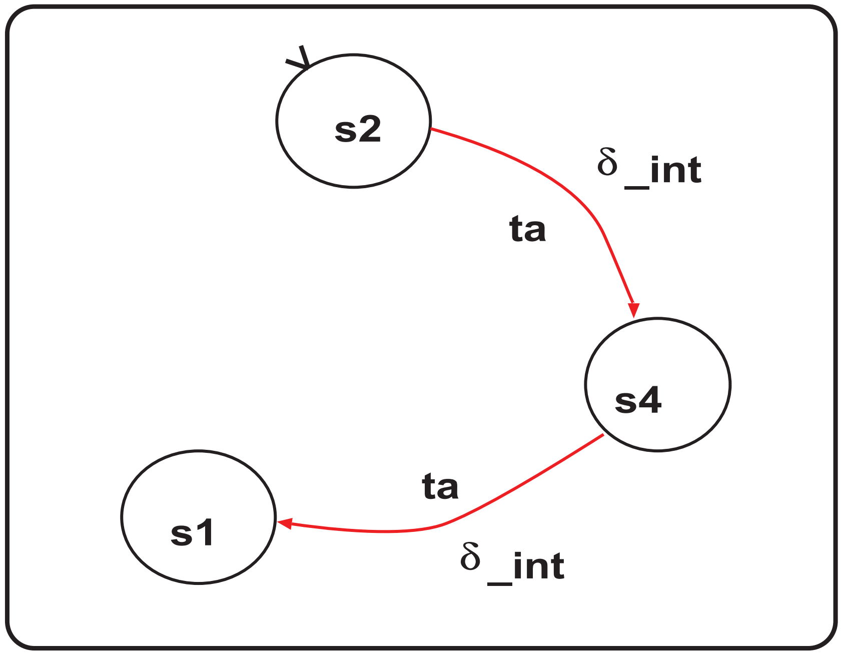
Event Graphs







DEVS without external events (model)



Operational Semantics

variables:
time = 0
current_state = s2

simulator:

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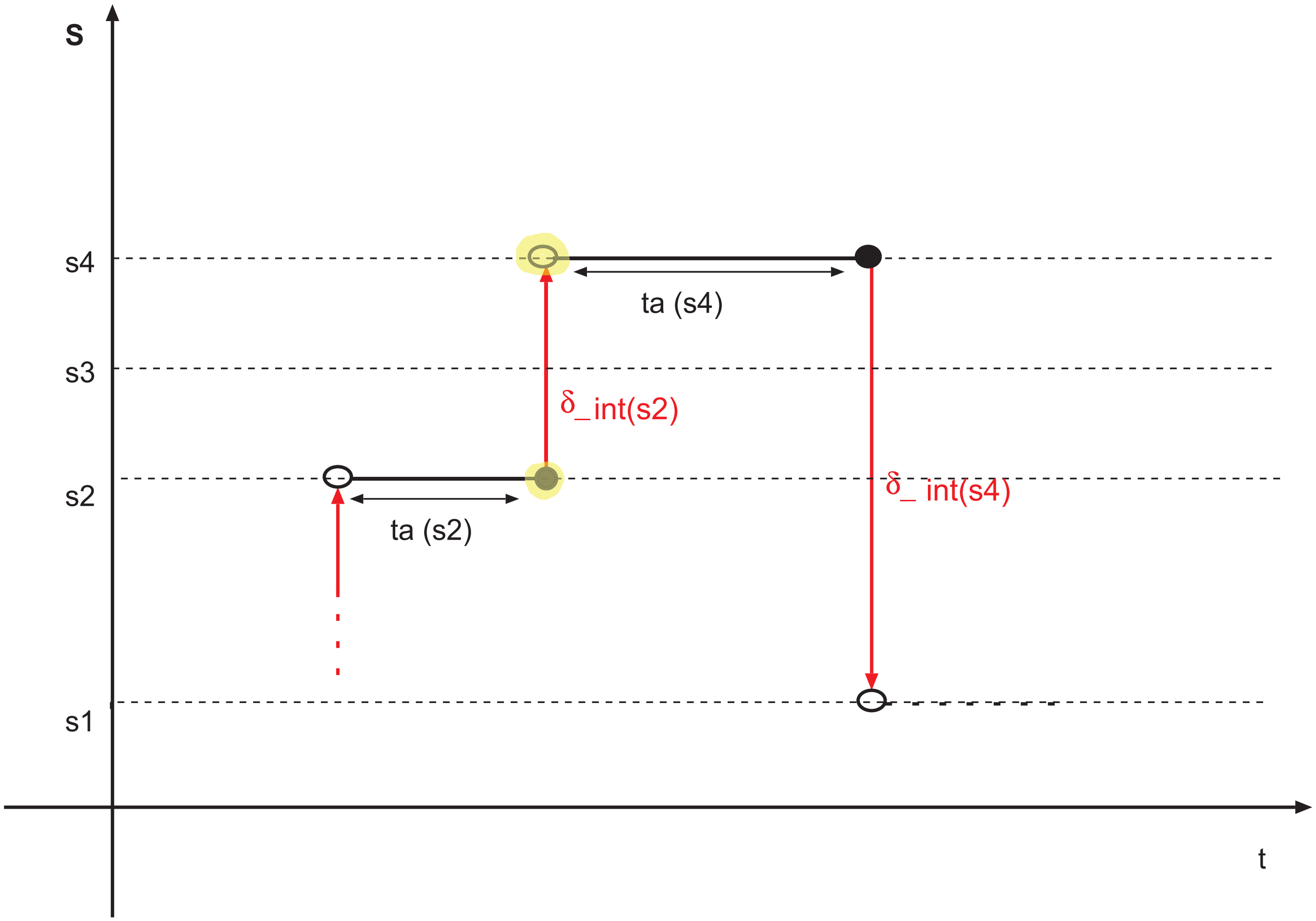
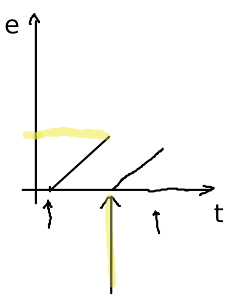
while True:
   time += ta(current_state)
   current_state =
            delta_int(current_state)
```

DEVS without external events (trajectories)



DEVS with external events (model)

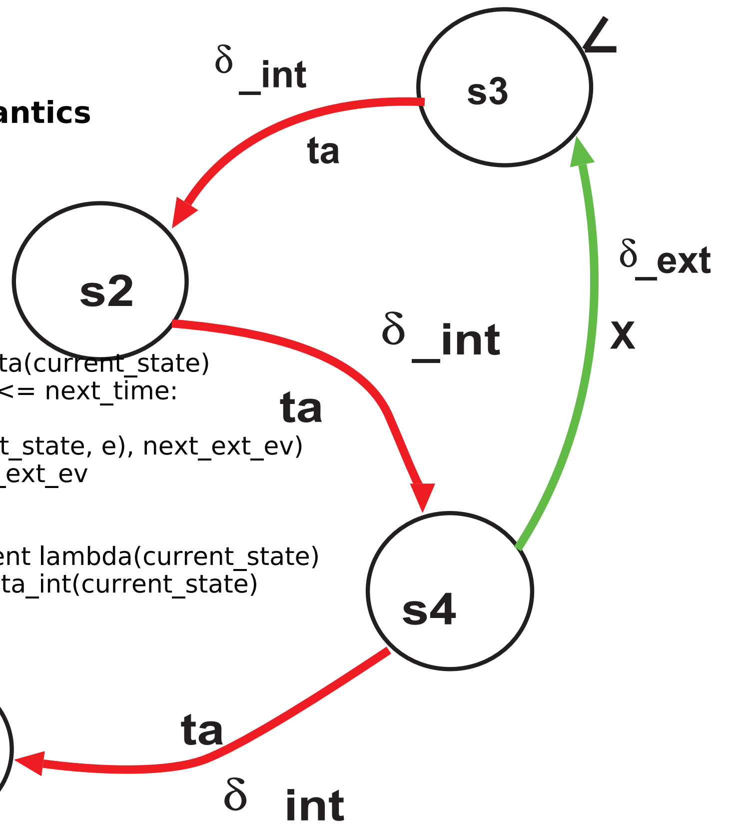


Operational Semantics

```
variables:
  time = 0
  current_state = s2

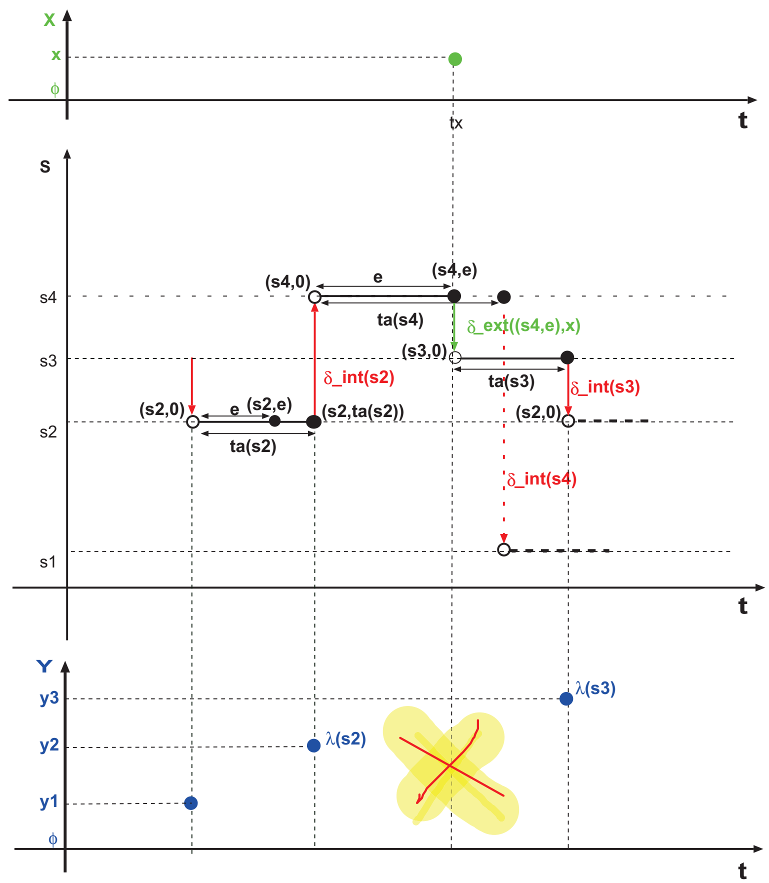
simulator:

  while True:
     next_time = time + ta(current_state)
     if time_next_ext_ev <= next_time:
        current_state =
          delta_ext((current_state, e), next_ext_ev)
         time = time_next_ext_ev
     else:
       time = next_time
       produce output event lambda(current_state)
       current_state = delta_int(current_state)
```

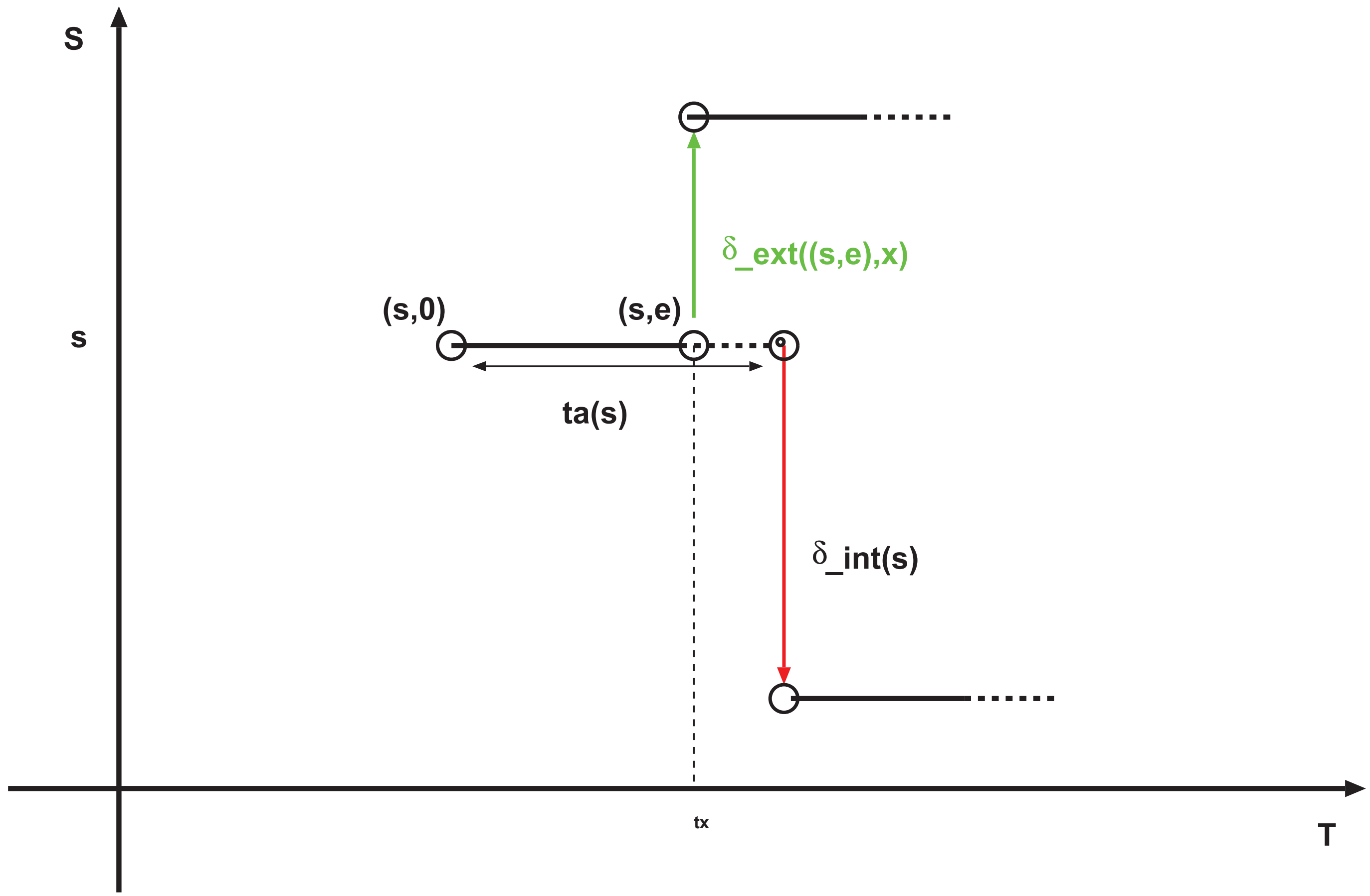


DEVS with external events (trajectories)

q= (s,e)



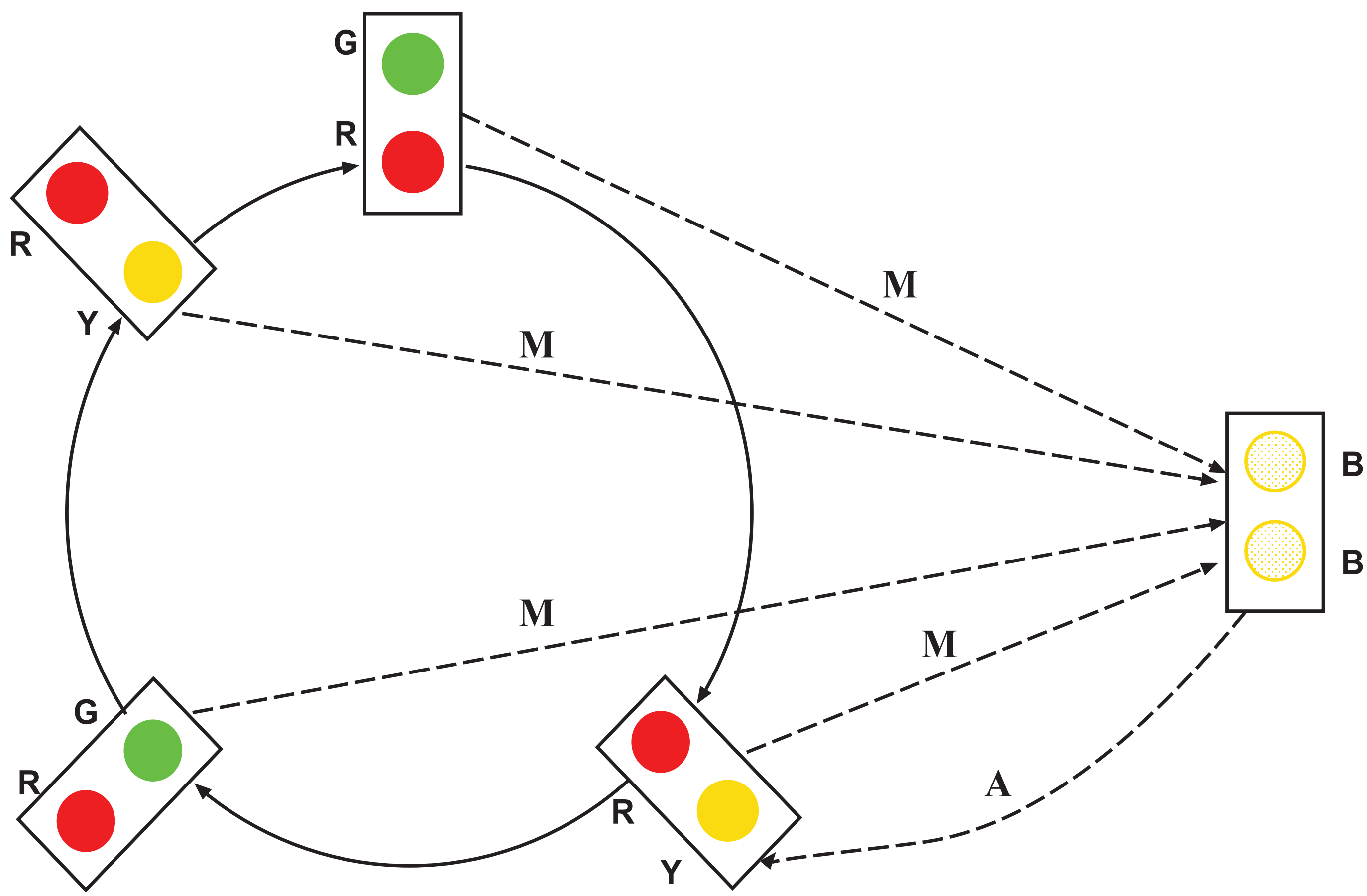
DEVS essence



$$DEVS = \langle X, S, Y, \delta_{int}, \delta_{ext}, \lambda, ta \rangle$$

$T = \mathbb{R}$	time base
X	input set
$\omega : T \rightarrow X \cup \{\phi\}$	input segment
S	state set
Y	output set
$\delta_{int} : S \rightarrow S$	internal transition function
$ta : S \rightarrow \mathbb{R}^{+}_{0,\infty}$	time advance function
$Q = \{(s, e) \mid s \in S, 0 \leq e \leq ta(s)\}$	total state, e is elapsed time
$\delta_{ext} : Q \times X \rightarrow S$	external transition function
$\lambda : S \rightarrow Y$	output function

Traffic Lights



$$trafficDEVS = < X, S, Y, \delta_{int}, \delta_{ext}, \lambda, ta >$$

$$T = \mathbb{R}$$

$$X = \{M, A\}$$

$$\omega : T \rightarrow X \cup \{\phi\}$$

$$S = \{RG, RY, GR, YR, BB\}$$

$$\delta_{int}(RG) = RY;\ \delta_{int}(RY) = GR$$

$$\delta_{int}(GR) = YR;\ \delta_{int}(YR) = RG$$

$$ta(RG) = 60s;\ ta(RY) = 10s$$

$$ta(GR) = 50s;\ ta(YR) = 10s$$

$$ta(BB) = +\infty$$

$$trafficDEVS = < X, S, Y, \delta_{int}, \delta_{ext}, \lambda, ta >$$

$$\delta_{ext}((RG, e), M) = BB$$

$$\delta_{ext}((RY, e), M) = BB$$

$$\delta_{ext}((GR, e), M) = BB$$

$$\delta_{ext}((YR, e), M) = BB$$

$$\delta_{ext}((BB, e), A) = RY$$

$$Y = \{GREY, YELLOW, BLINK\}$$

$$\lambda(RG) = \lambda(RY) = \lambda(GR) = GREY$$

$$\lambda(YR) = YELLOW$$

$$\lambda(BB) = BLINK$$

Coupled DEVS

$$coupledDEVS \equiv \langle X_{self}, Y_{self}, D, \{M_i\}, \{I_i\}, \{Z_{i,j}\}, select \rangle$$

$$\{M_i | i \in D\}.$$

$$M_i = \langle S_i, ta_i, \delta_{int,i}, X_i, \delta_{ext,i}, Y_i, \lambda_i \rangle, \forall i \in D.$$

$$\{I_i | i \in D \cup \{self\}\}.$$

$$\forall i \in D \cup \{self\} : I_i \subseteq D \cup \{self\}.$$

$$\forall i \in D \cup \{self\} : i \notin I_i.$$

I_i are the influenc**ee** sets describing the connection topology

$Z_{i,j}$ output-to-input translation

$$\{Z_{i,j} | i \in D \cup \{self\}, j \in I_i\},$$

$$\begin{aligned}
Z_{self,j} &: X_{self} \rightarrow X_j &&, \forall j \in D, \\
Z_{i,self} &: Y_i \rightarrow Y_{self} &&, \forall i \in D, \\
Z_{i,j} &: Y_i \rightarrow X_j &&, \forall i, j \in D.
\end{aligned}$$

Together, I_i and $Z_{i,j}$ completely specify the coupling (structure and behaviour)

Tie-breaking among simultaneous events

$$select : 2^D \rightarrow D$$

Choose a unique component from any non-empty subset E of D:

$$select(E) \in E.$$

E corresponds to the set of all components having a state transition simultaneously (*collisions*).

Closure under coupling

From the coupled DEVS

$$\langle X_{self}, Y_{self}, D, \{M_i\}, \{I_i\}, \{Z_{i,j}\}, select\rangle,$$

with all components M_i atomic DEVS models

$$M_i = \langle S_i, ta_i, \delta_{int,i}, X_i, \delta_{ext,i}, Y_i, \lambda_i\rangle, \forall i \in D$$

the atomic DEVS

$$\langle S, ta, \delta_{int}, X, \delta_{ext}, Y, \lambda\rangle$$

is constructed.

Closure: state and time-advance

$$S = \times_{i \in D} Q_i,$$

where

$$Q_i = \{(s_i, e_i) | s \in S_i, 0 \leq e_i \leq ta_i(s_i)\}, \forall i \in D.$$

$$ta : S \rightarrow \mathbb{R}^+_{0,+\infty}$$

Select the *most imminent* event time, = smallest time *remaining* until internal transition, of all the components

$$ta(s) = min\{\sigma_i = ta_i(s_i) - e_i | i \in D\}.$$

Dealing with simultaneous events

Imminent components:

$$IMM(s) = \{i \in D | \sigma_i = ta(s)\}.$$

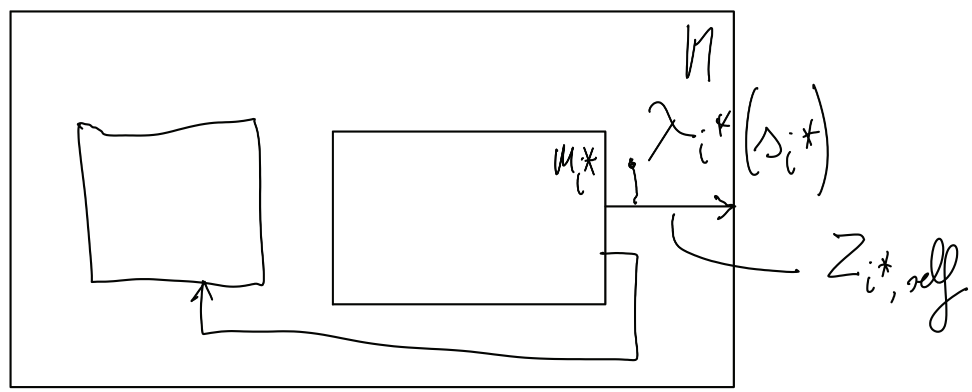
$select$ *one* component i^* of the coupled model

$$\begin{array}{rccl} select & : & 2^D & \rightarrow & D \\ & & IMM(s) & \rightarrow & i^* \end{array}$$

Output (at internal transition time)

$$\lambda(s) = \begin{cases} Z_{i^*,self}(\lambda_{i^*}(s_{i^*})) & \text{,if } self \in I_{i^*}, \\ \phi & \text{,if } self \notin I_{i^*}. \end{cases}$$

Conceptually, the non-event ϕ is generated if i^* is not connected to the output of the coupled model.



Internal transition function

$$\overbrace{((s_1,e_1),(s_2,e_2),\ldots,(s_n,e_n))}^{s}$$

$$\delta_{int}(s) = (\ldots,(s'_j,e'_j),\ldots), \text{ where}$$

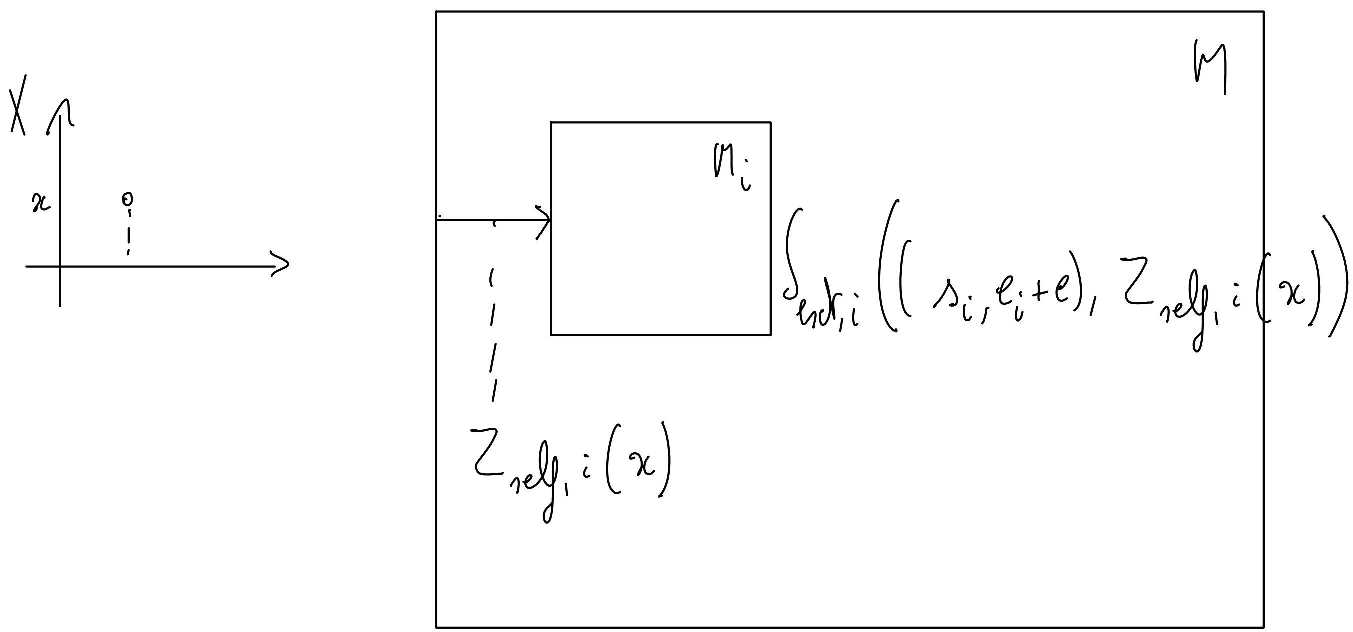
$$
\begin{aligned}
(s'_j,e'_j) &= (\delta_{int,j}(s_j),0) && \text{, for } j = i^*, \\
&= (\delta_{ext,j}(s_j, e_j + ta(s), Z_{i^*,j}(\lambda_{i^*}(s_{i^*}))),0) && \text{, for } j \in I_{i^*} \\
& && \text{(and } Z_{i^*,j}(\lambda_{i^*}(s_{i^*})) \neq \\
&= (s_j, e_j + ta(s)) && \text{, otherwise.}
\end{aligned}
$$

External transition function

$$\delta_{ext}(s, e, x) = (\ldots, (s'_i, e'_i), \ldots), \text{ where}$$

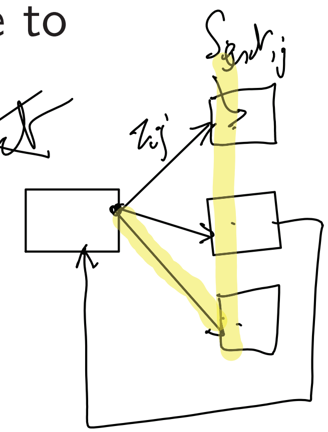
$$\begin{aligned}(s'_i, e'_i) &= (\delta_{ext,i}(s_i, \underline{e_i + e}, Z_{self,i}(x)), 0) &&, \text{for } i \in I_{self}, \\ &= (s_i, \underline{e_i + e}) &&, \text{otherwise.}\end{aligned}$$

$i \in I_{self}$



DEVS limitations

- a conflict due to simultaneous internal and external events is resolved by ignoring the internal event. It should be possible to explicitly specify behaviour in case of conflicts;
- there is limited potential for parallel implementation;
- the *select* function is an artificial legacy of the semantics of traditional sequential simulators based on an event list;
- it is not possible to explicitly describe variable structure.



DS-DEVS

Some of these are resolved in *parallel DEVS*

DEVS Solver

- Iterative simulation of DEVS model
- Possibly distributed implementation

message m	**simulator**	**coordinator**
$(*, from, t)$	simulator correct only if $t = t_N$	
	$y \leftarrow \lambda(s)$	**send** $(*, self, t)$ to i^*, where
	if $y \neq \phi$:	$i^* = select(imm\_children)$
	send $(\lambda(s), self, t)$ to $parent$	$imm\_children = \{i \in D \| M_i.t_N = t\}$
	$s \leftarrow \delta_{int}(s)$	$active\_children \leftarrow active\_children \cup \{i^*\}$
	$t_L \leftarrow t$	
	$t_N \leftarrow t_L + ta(s)$	
	send $(done, self, t_N)$ to $parent$	

message m	simulator	coordinator
$(x, from, t)$	simulator correct only if $t_L \leq t \leq t_N$ (ignore δ_{int} to resolve a $t = t_N$ *conflict*)	
	$e \leftarrow t - t_L$ $s \leftarrow \delta_{ext}(s, e, x)$ $t_L \leftarrow t$ $t_N \leftarrow t_L + ta(s)$ **send** $(done, self, t_N)$ to $parent$	$\forall i \in I_{self}$: **send** $(Z_{self,i}(x), self, t)$ to i $active\_children \leftarrow active\_children \cup \{i$

message m	simulator	coordinator
$(y, from, t)$		$\forall i \in I_{from} \setminus \{self\}$: **send** $(Z_{from,i}(y), from, t)$ to i $active\_children \leftarrow active\_children \cup \{i\}$ **if** $self \in I_{from}$: **send** $(Z_{from,self}(y), self, t)$ to $parent$
$(done, from, t)$		$active\_children \leftarrow active\_children \setminus \{from\}$ **if** $active\_children = \emptyset$: $t_L \leftarrow t$ $t_N \leftarrow min\{M_i.t_N \mid i \in D\}$ **send** $(done, self, t_N)$ to $parent$

DEVS simulator main loop

$t \leftarrow t_N$ of topmost coordinator

repeat until $t \geq t_{end}$ (or some other termination condition)

  send $(*, main, t)$ to topmost coupled model top

  wait for $(done, top, t_N)$

  $t \leftarrow t_N$
