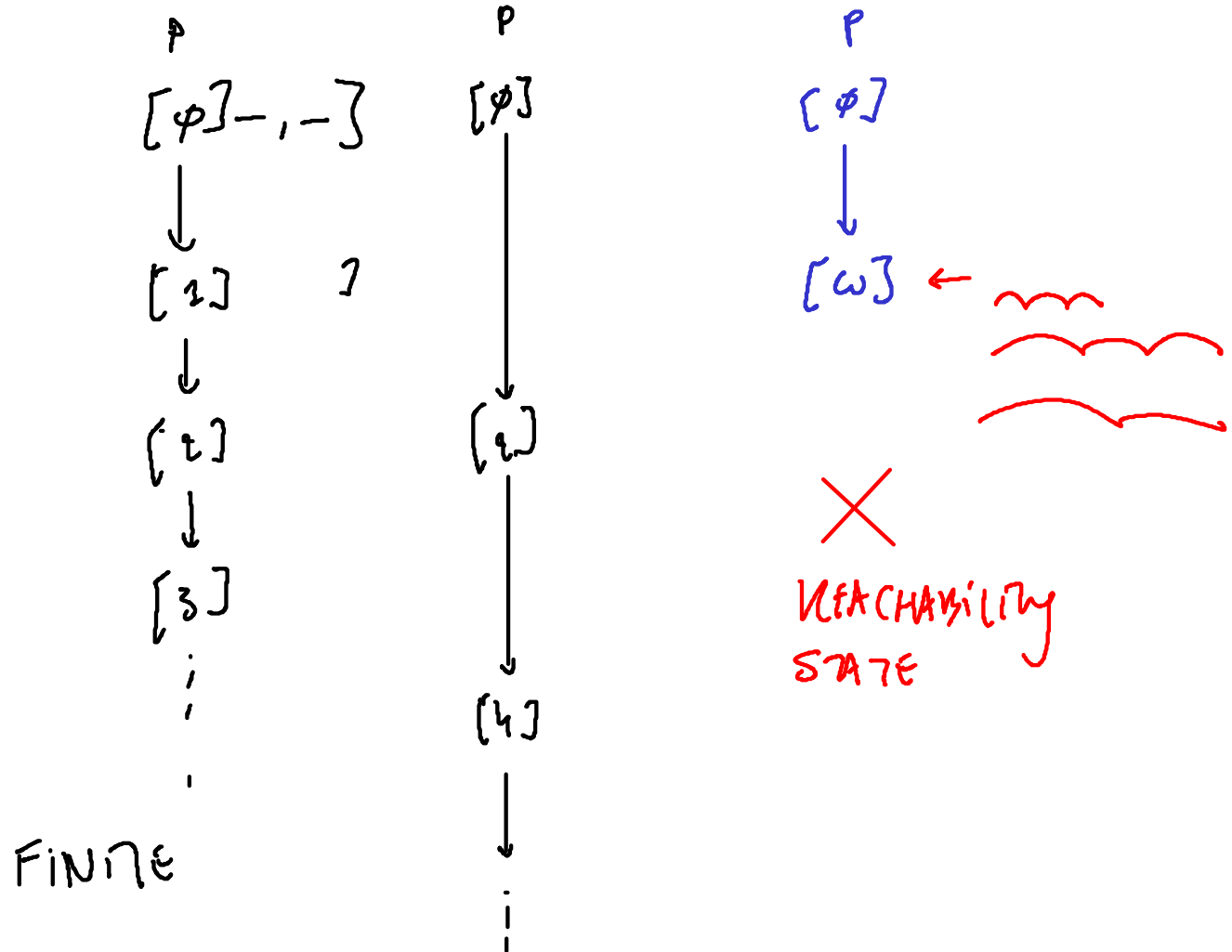


ABSTRACTION

reachability

coverability



FINITE

REACHABLE ($[2]$) ?
 $[3, \omega, \ast]$

$$\begin{array}{cccc}
 P_0 & P_1 & P_2 & P_3 \\
 [1, & 0, & 0, & 0]
 \end{array}$$

$$\downarrow T_\phi$$

$$\rightarrow [0, 1, 1, 0]$$

$$\begin{array}{l}
 \swarrow T_1 \\
 \searrow T_2
 \end{array}$$

$$[0, \cancel{0}, 1, 1]$$

$$\uparrow T_1 \rightarrow [1, 0, \omega, 0]$$

$$\downarrow T_\phi$$

$$\rightarrow [\cancel{0}, 1, \omega, 0]$$

$$\downarrow T_1$$

$$\searrow T_2$$

$$[\cancel{1}, \cancel{0}, \omega, 0]$$

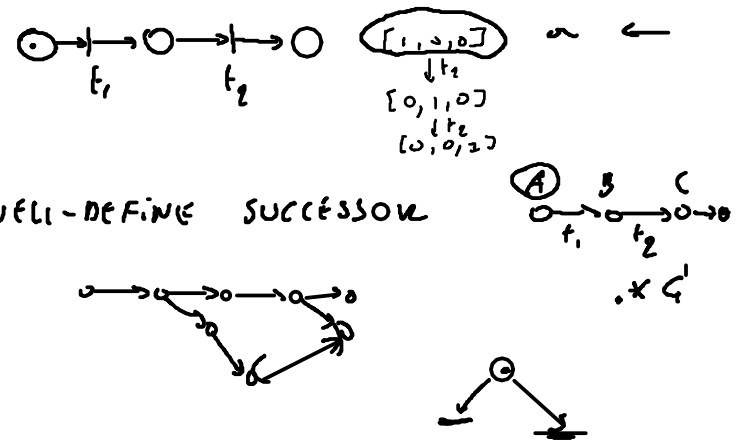
$$[0, 0, \omega, 1]$$

- "TEMPORAL" LOGIC
- LINEAR LOGIC
- BRANCHING LOGIC

ORDERING OF EVENTS

EACH MOMENT IN "TIME" HAS A WELL-DEFINE SUCCESSOR

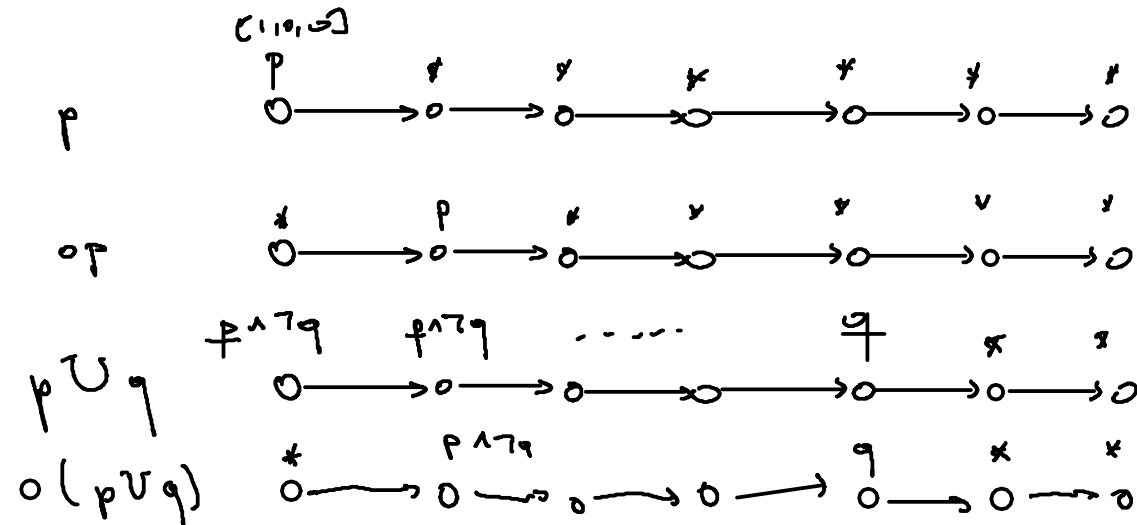
MULTIPLE POSSIBLE TIME COURSES



LTL: LINEAR TEMPORAL LOGIC (PNUFF 1970s)

CTL: COMBINATION TREE LOGIC (BRANCHING TIME)

PROPOSITION $\varphi ::= \text{TRUE} \mid \text{atomic } p \mid \varphi_1 \wedge \varphi_2 \mid \neg \varphi \mid \text{NEXT } \varphi \mid \text{UNTIL } \varphi_1 \cup \varphi_2$

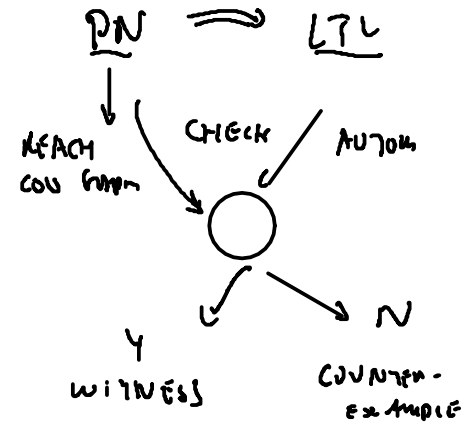


NEXT UNTIL
 φ $\varphi_1 \cup \varphi_2$
 X

LTL

SATISF?

ooo p

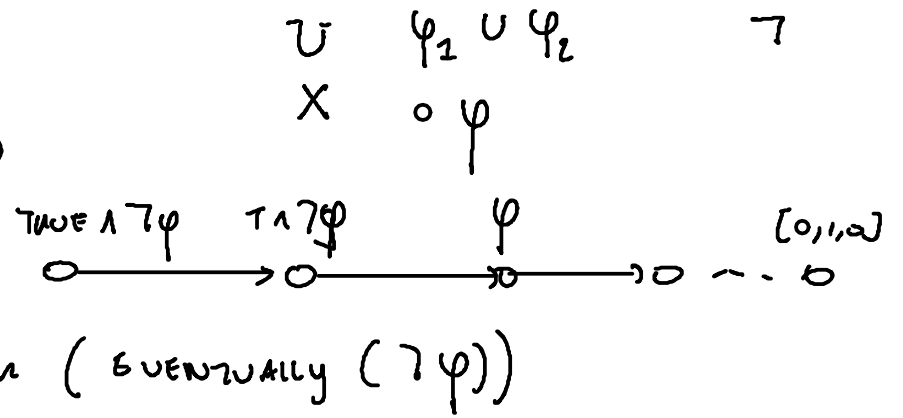


$\neg \varphi \cup \varphi$
 $\neg \varphi \wedge \neg \varphi$

DEFINING OPERATIONS

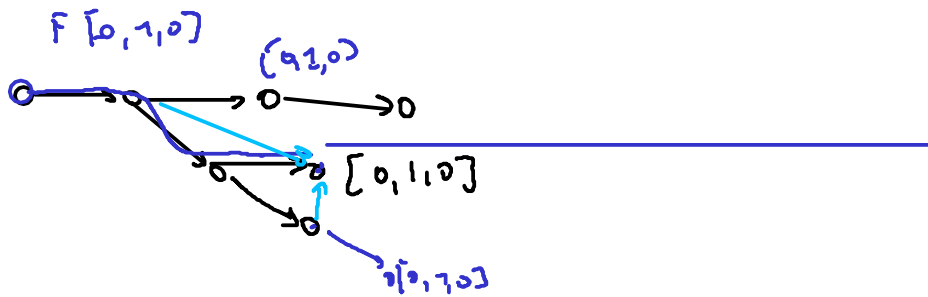
EVENTUALLY ψ $F \Diamond \psi := \text{TRUE} \cup \psi$

ALWAYS (GLOBALLY) ψ $G \Box \psi :=$



$F [0,1,0]$

$\neg \Diamond (\neg \psi)$
|
NOT



\forall PATHS $A F [0,2,0]$
 \exists A PATH $E F [0,7,0]$

(TLX)

