

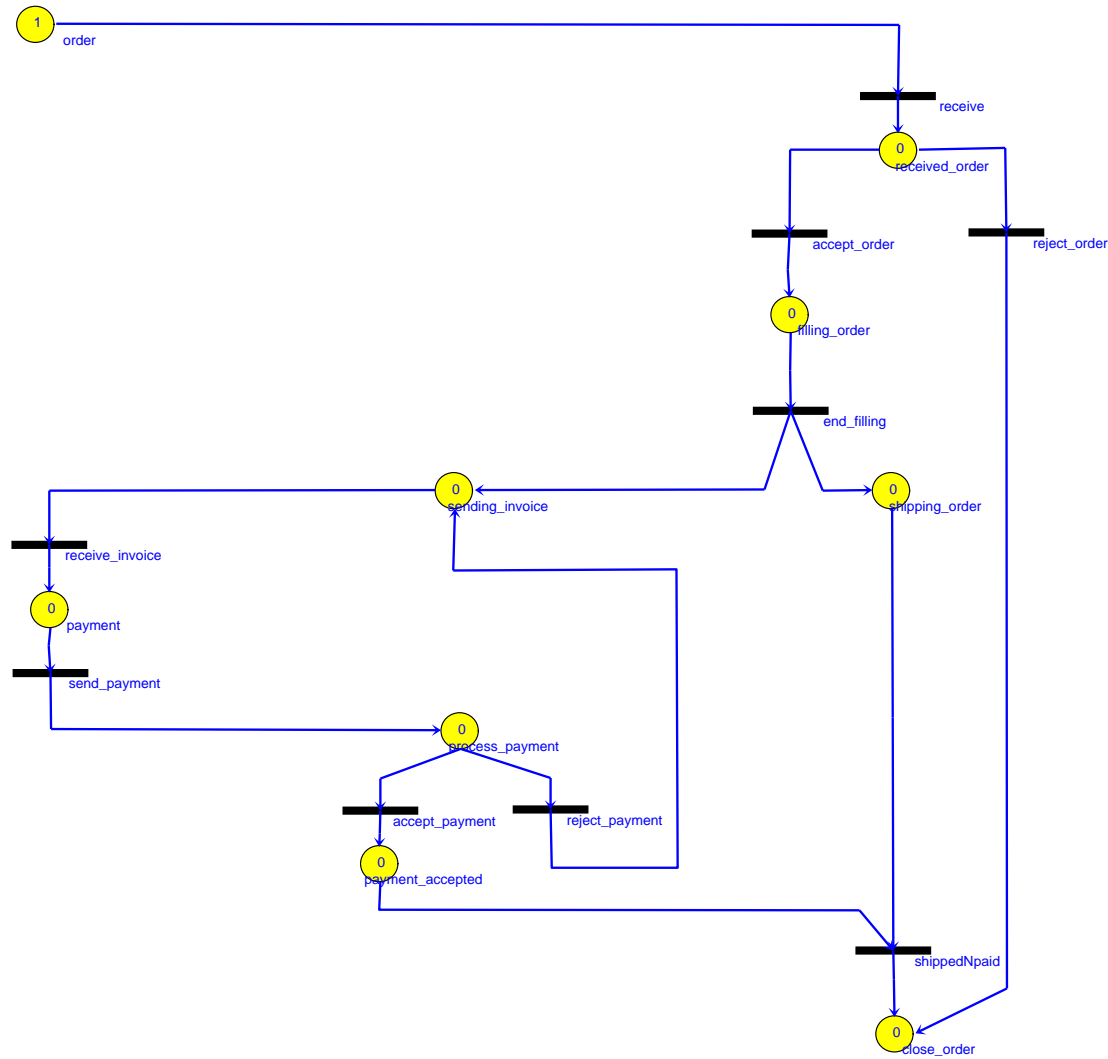
# Activity Diagrams

- describe *behaviour*
- at *high level of abstraction*
- focus on *workflows (processes/activities)*
- elegant description of *concurrency*
- can express *non-determinism*
- as of UML 2.0 based on *Petri Nets* (before: Statecharts)

# Petri nets

- C.A. Petri, Kommunikation mit Automaten, Ph.D. Thesis, Schriften des Institutes für Instrumentelle Mathematik, Bonn, 1962.
- Formalism based on FSA
- Graphical notation
- Additions to FSA:
  - Explicitly (graphically) represent when event is enabled  
→ describe control logic
  - Elegant notation for concurrency
  - Elegant notation for synchronization
  - Express non-determinism

# Order Processing Example

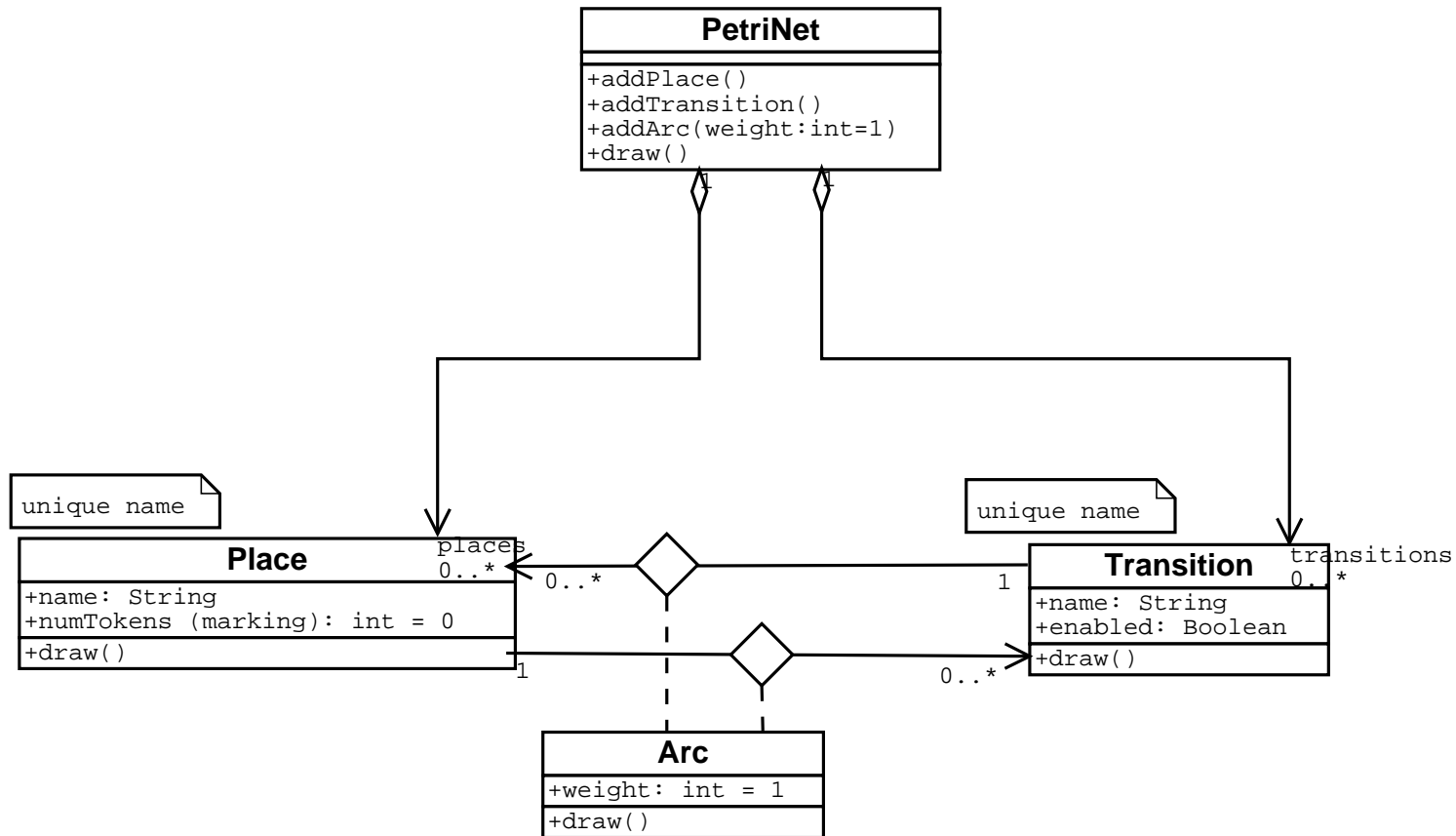


# Petri net notation and definition (no dynamics)

$$(P, T, A, w)$$

- $P = \{p_1, p_2, \dots\}$  is a finite set of *places*
- $T = \{t_1, t_2, \dots\}$  is a finite set of *transitions*
- $A \subseteq (P \times T) \cup (T \times P)$  is a set of *arcs*
- $w : A \rightarrow \mathbb{N}$  is a *weight function*

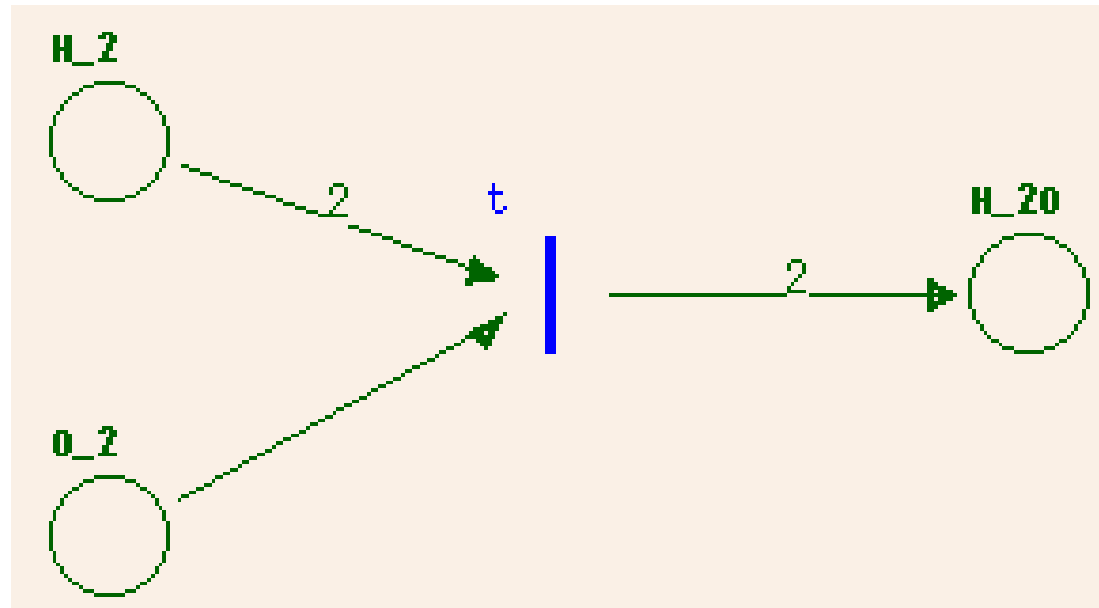
# Class Diagram model of Petri Net Abstract Syntax



# Derived Entities

- $I(t_j) = \{p_i : (p_i, t_j) \in A\}$  set of *input places* to transition  $t_j$   
( $\equiv$  conditions for transition)
- $O(t_j) = \{p_i : (t_j, p_i) \in A\}$  set of *output places* from transition  $t_j$   
( $\equiv$  affected by transition)
- Transitions  $\equiv$  events
- similarly: input- and output-transitions for  $p_i$
- graphical representation (concrete syntax):  
*Petri net graph* (multigraph)

# Example Petri net



- $P = \{H_2, O_2, H_2O\}$
- $T = \{t\}$
- $A = \{(H_2, t), (O_2, t), (t, H_2O)\}$
- $w((H_2, t)) = 2, w((O_2, t)) = 1, w((t, H_2O)) = 2$

# Introducing State: Petri net Markings

- Conditions met ? Use *tokens* in places
- Token assignment  $\equiv$  *marking*  $x$

$$x : P \rightarrow \mathbb{N}$$

- A marked Petri net

$$(P, T, A, w, x_0)$$

$x_0$  is the *initial marking*

- The *state*  $\mathbf{x}$  of a marked Petri net

$$\mathbf{x} = [x(p_1), x(p_2), \dots, x(p_n)]$$

Number of tokens need not be bounded (cfr. State Automata states).



# State Space of Marked Petri net

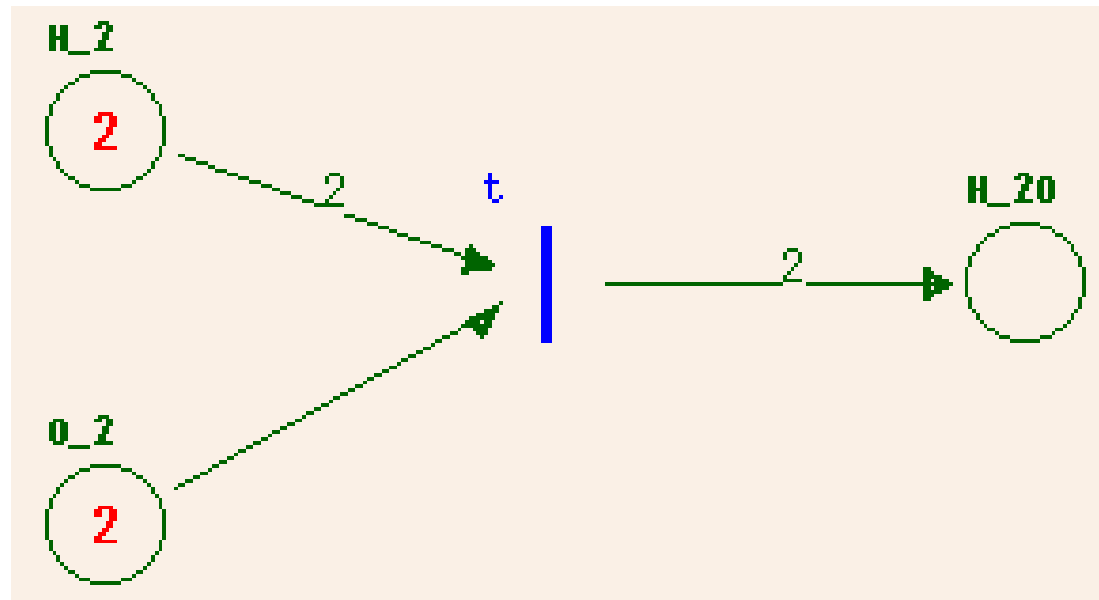
- All  $n$ -dimensional vectors of nonnegative integer markings

$$X = \mathbb{N}^n$$

- Transition  $t_j \in T$  is *enabled* if

$$x(p_i) \geq w(p_i, t_j), \forall p_i \in I(t_j)$$

# Example with marking, enabled



# Petri Net Dynamics

State Transition Function  $f$  of marked Petri net  $(P, T, A, w, x_0)$

$$f : \mathbb{N}^n \times T \rightarrow \mathbb{N}^n$$

is defined for transition  $t_j \in T$  if and only if

$$x(p_i) \geq w(p_i, t_j), \forall p_i \in I(t_j)$$

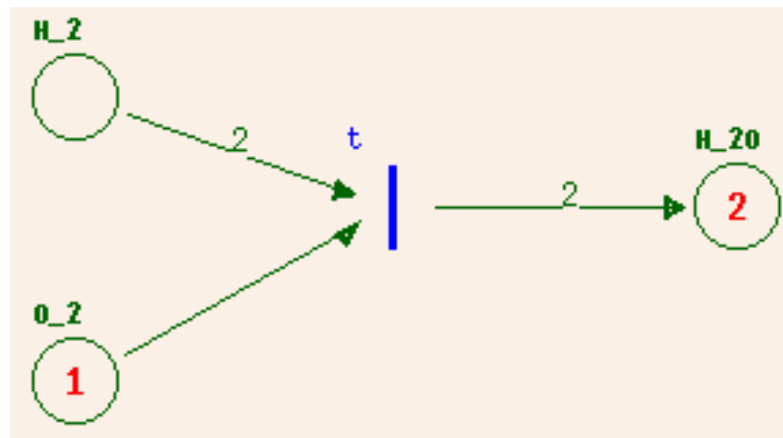
If  $f(\mathbf{x}, t_j)$  is defined, set  $\mathbf{x}' = f(\mathbf{x}, t_j)$  where

$$x'(p_i) = x(p_i) - w(p_i, t_j) + w(t_j, p_i)$$

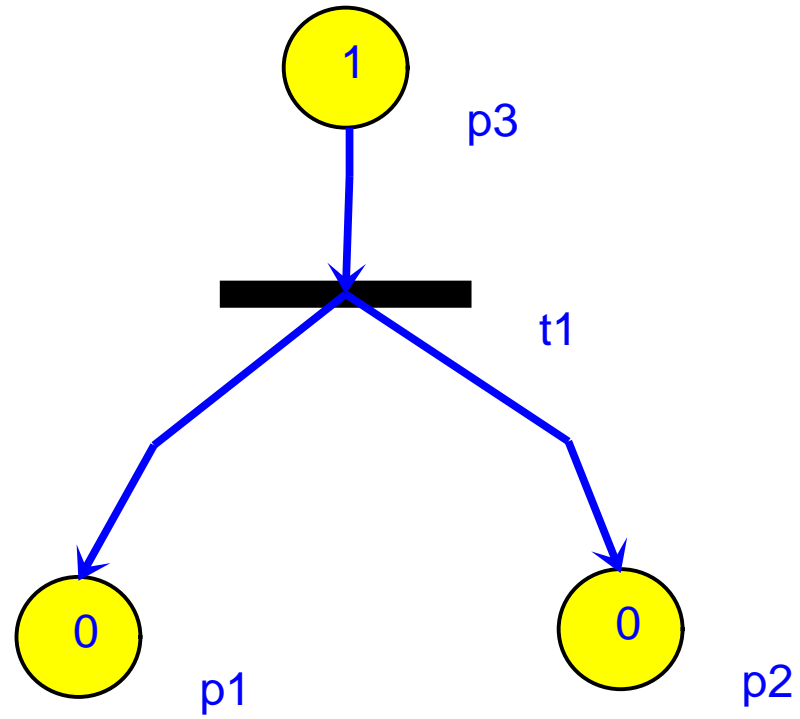
- State transition function  $f$  based on *structure* of Petri net
- Number of tokens *need not be conserved* (but can)

# Example “firing”

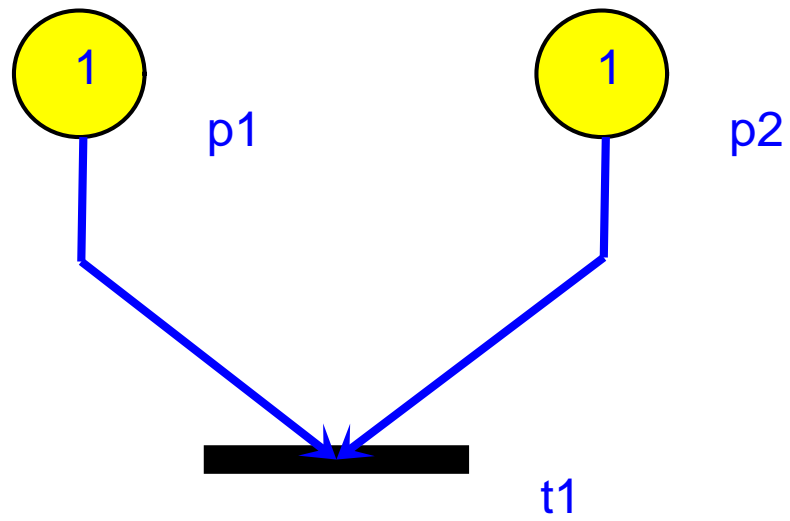
- Sequential Manual execution
- Transition:  $[2, 2, 0] \rightarrow [0, 1, 2]$



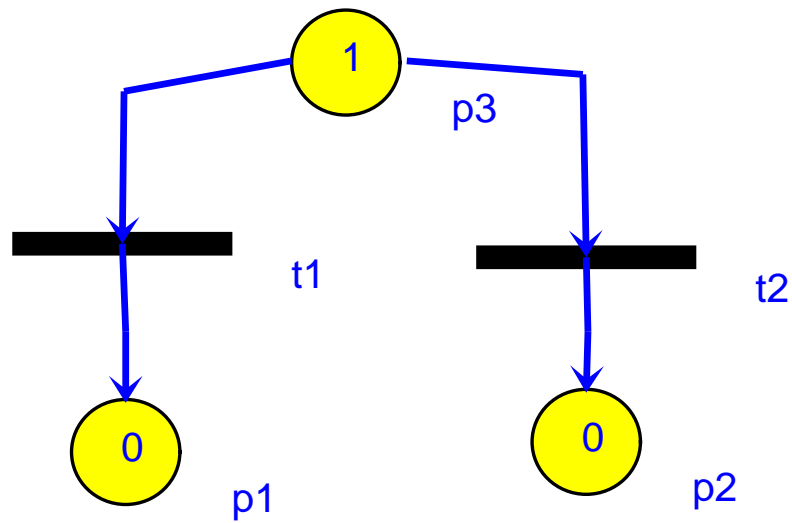
# Fork



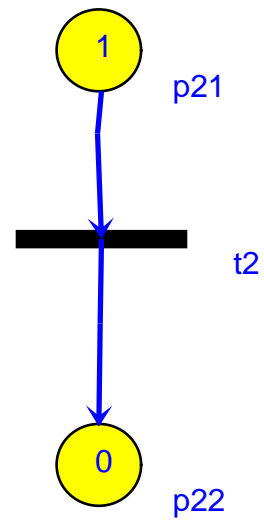
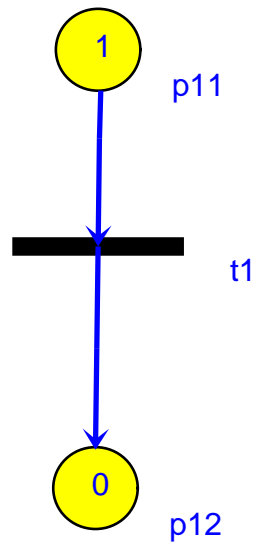
# Join



# Conflict, choice, decision



# Concurrency

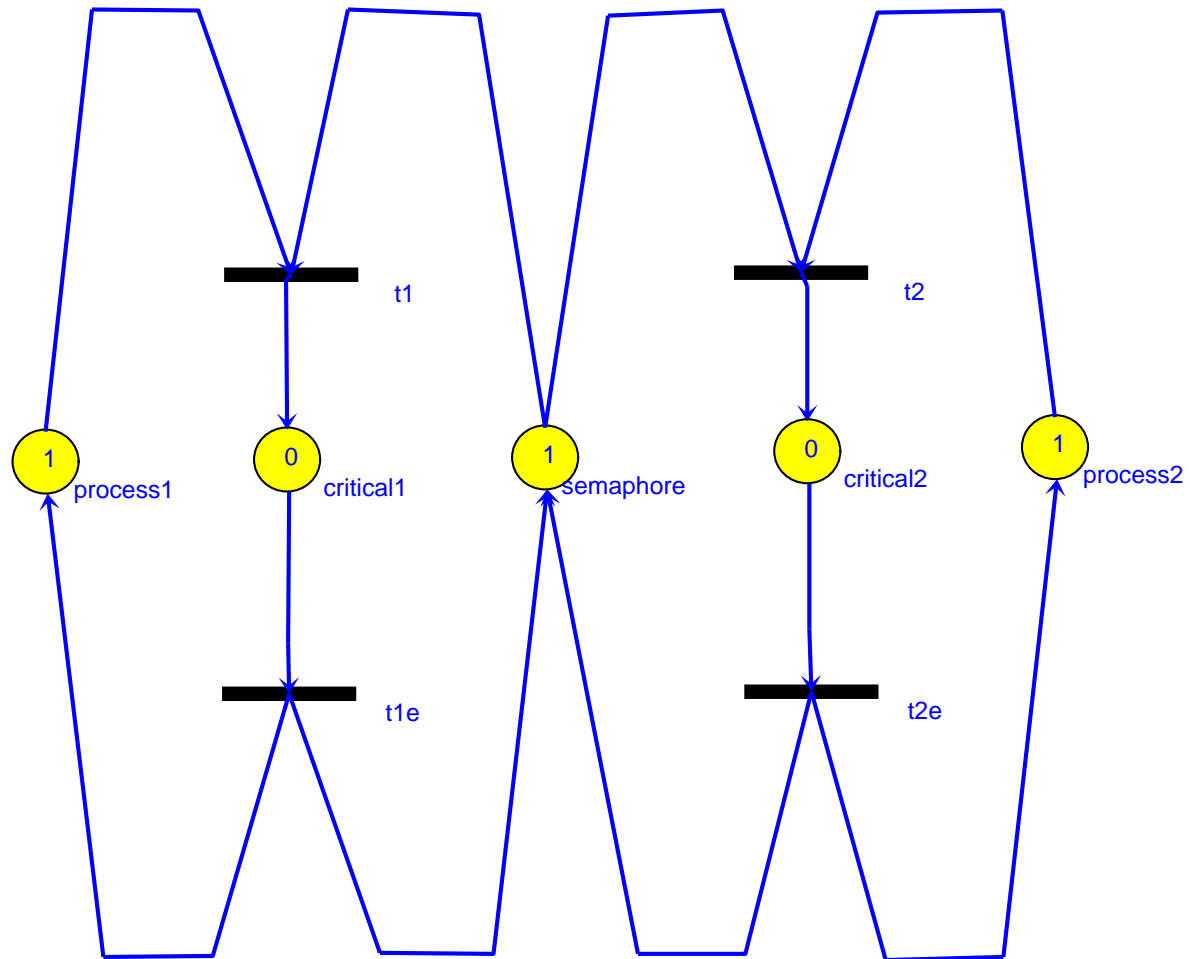




# Semantics

- *sequential vs. parallel*
- Handle nondeterminism:
  1. User choice
  2. Priorities
  3. Probabilities (Monte Carlo)
  4. Reachability Graph (enumerate all choices)

# Application: Critical Section



# Reachability Graph

